

Ways of measuring errors: $x \in \mathbb{R}$ is the exact value and x^* the approximation

① absolute error $e_{abs} = |x - x^*|$

② relative error ... $e_{rel} = \frac{e_{abs}}{|x|}$

Types:

approximation	approximand	type of error
v^*	v	initial error
$M(v^*)$	$M(v)$	propagated error
$M^*(v^*)$	$M^*(v^*)$	generated error
$M^*(v^*)$	$M(v)$	total cumulative error

what we get what we want

M is the function or algorithm we want to apply.

M^* is an approximation of that function programmed into a computer..

already initial error in the inputs
 ↓
 add the exponents

```
julia> 4.27*3.68
15.7136
```

$$(4.27 \times 10^1) \times (3.68 \times 10^1) = 15.7136 \times 10^2 \rightarrow 1.57 \times 10^3$$

rounding step 35 generates more error..

```
julia> 4.27e1*3.68e1
1571.36
```

round back to 35
 given $1570. = 1.57 \times 10^3$

Because of the rounding the computer doesn't perform multiplication exactly...and so ...

Addition of a small to a large # may have no effect..

$$5.18 \times 10^2 + 4.37 \times 10^{-1} = 5.18 \times 10^2 + 0.00437 \times 10^2$$

$$= 5.18437 \times 10^2 \rightarrow 5.18 \times 10^2$$

How much smaller does the have to be for no effect?


```
julia> epsilon=1e-16
1.0e-16

julia> 1+epsilon==1
true

julia> epsilon=1e-15
1.0e-15

julia> 1+epsilon==1
false
```

← too small

← makes a difference ..

1.0
0.0000000000000001

indicates about 15 significant digits are available ..

also :

the associative law for addition does not always hold.

$$(a+b)+c = a+(b+c)$$

If approximate addition... it can happen that

$$(a+b)+c \neq a+(b+c)$$

```
julia> a=1.0
1.0

julia> b=2.0
2.0

julia> c=3.0
3.0

julia> (a+b)+c==a+(b+c)
true
```

← In this case associativity holds...

```

julia> a=0.1
0.1
julia> b=0.2
0.2
julia> c=0.3
0.3
julia> (a+b)+c==a+(b+c)
false

```

there is actually initial error here...

$a \approx 0.1$ but not exact because working in base 2...

Non associative, because of the rounding errors...

```

julia> (a+b)+c
0.60000000000000001
julia> a+(b+c)
0.6

```

this looks strange because the computer is working base 2 behind the scenes...

Scientific notation base 10,

$$.1 = 1.0 \times 10^{-1}$$

if this were

$$.1 = m \times 2^n$$

never get a factor of 10 from powers of 2...

this simple number in decimal has a repeating binary expansion...

any fraction that doesn't have a power of 2 in the denominator turns into a repeating binary expansion.

recall for decimal numbers, fractions of the form

$$\frac{p}{2^n 5^m} \quad \text{all terminate}$$

but not the others...

- [Step 01](#)
- [Step 02](#)
- [Step 03, Non-Associativity of Addition, Types of Errors](#)
- [Step 04, Realistic Floating Point Data Types, Rounding Examples](#)
- [Step 05, Taylor's Theorem](#)

↑ need to have for wednesday...

already initial error in the inputs
 ↓
 add the exponents

$$\underbrace{(4.27 \times 10^1)}_{a^*} \times \underbrace{(3.68 \times 10^1)}_{b^*} = 15.7136 \times 10^2 \rightarrow 1.57 \times 10^3$$

↑ generated...

```
julia> 4.27*3.68
15.7136
```

$$e_{rel}(a) \leq 5 \times 10^{-3}$$

$$e_{rel}(b) \leq 5 \times 10^{-3}$$

→ propagated error is sum of rel. error in multiplication

$$e_{rel}(a+b) \leq 10 \times 10^{-3} = 1 \times 10^{-2}$$