

Recall: Solving $y' = f(t, y)$ such that $y(t_0) = y_0$.

Started with $O(h^2)$ for how much error is generated at each step...

since $h = \frac{T-t_0}{N}$
 then $N = \frac{T-t_0}{h}$
 and after N intervals one of the h 's is gone...

$$\|e_{n+1}\| \leq (1+h\lambda)\|e_n\| + O(h^2)$$

Thus

$$\|e_{n+1}\| \leq (1+h\lambda)\|e_n\| + ch^2$$

Use induction...

$e_0 = y_0 - y(t_0) \approx 0$ up to rounding errors..

$$\|e_1\| \leq (1+h\lambda)\|e_0\| + ch^2$$

$$\|e_2\| \leq (1+h\lambda)\|e_1\| + ch^2$$

$$\leq (1+h\lambda) \left\{ (1+h\lambda)\|e_0\| + ch^2 \right\} + ch^2$$

$$= (1+h\lambda)^2 \|e_0\| + \{1 + (1+h\lambda)\} ch^2$$

$$\|e_3\| \leq (1+h\lambda)\|e_2\| + ch^2$$

$$\leq (1+h\lambda) \left\{ (1+h\lambda)^2 \|e_0\| + \{1 + (1+h\lambda)\} ch^2 \right\} + ch^2$$

$$= (1+h\lambda)^3 \|e_0\| + \underbrace{(1 + (1+h\lambda) + (1+h\lambda)^2)}_{\text{geometric sum}} ch^2$$

one less than 3

In general

$$\|e_n\| \leq (1+h\lambda)^n \|e_0\| + \sum_{k=0}^{n-1} (1+h\lambda)^k ch^2$$

simplify this geometric sum...

$$S_n = \sum_{k=0}^{n-1} \alpha^k \quad \leftarrow \text{simplify this...}$$

$$\alpha S_n = \alpha + \alpha^2 + \alpha^3 + \dots + \alpha^n$$

$$S_n = 1 + \alpha + \alpha^2 + \dots + \alpha^{n-1}$$

$$(\alpha - 1) S_n = \alpha^n - 1$$

Thus $\sum_{k=0}^{n-1} \alpha^k = \frac{\alpha^n - 1}{\alpha - 1}$

Therefore

$$\|e_n\| \leq (1+h\lambda)^n \|e_0\| + \sum_{k=0}^{n-1} (1+h\lambda)^k ch^2$$

$$= (1+h\lambda)^n \|e_0\| + \frac{(1+h\lambda)^n - 1}{(1+h\lambda) - 1} ch^2$$

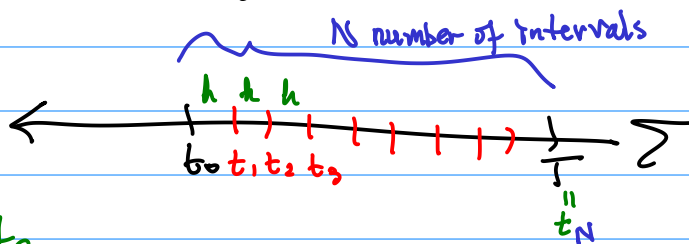
$$= (1+h\lambda)^n \|e_0\| + \frac{(1+h\lambda)^n - 1}{\lambda} ch$$

Now I want to work with $(1+h\lambda)^n$ and bound this...

Remember the setup

$$y' = f(t, y) \quad y(t_0) = y_0$$

Made a grid. Solving on $[t_0, T]$ and



$$h = \frac{T - t_0}{N}$$

$$t_n = t_0 + hn$$

$$hn = t_n - t_0$$

Check that $t_N = T$

$$t_N = t_0 + hN = t_0 + \frac{T - t_0}{N} N = T$$

$$\|e_n\| \leq (1+h\lambda)^n \|e_0\| + \frac{(1+h\lambda)^n - 1}{\lambda} ch$$

$$e^{h\lambda} = 1 + h\lambda + \frac{1}{2}(h\lambda)^2 + \frac{1}{3!}(h\lambda)^3 + \frac{1}{4!}(h\lambda)^4 + \dots$$

all these terms are positive since $h\lambda > 0$

Thus,

$$1 + h\lambda \leq e^{h\lambda}$$

$$\|e_n\| \leq (e^{h\lambda})^n \|e_0\| + \frac{ch}{\lambda} [(e^{h\lambda})^n - 1]$$

$$= e^{hn\lambda} \|e_0\| + \frac{ch}{\lambda} [e^{hn\lambda} - 1]$$

$$= e^{(t_n - t_0)\lambda} \|e_0\| + \frac{ch}{\lambda} [e^{(t_n - t_0)\lambda} - 1]$$

Then

$$\|e_n\| \leq e^{(T - t_0)\lambda} \|e_0\| + \frac{ch}{\lambda} [e^{(T - t_0)\lambda} - 1]$$

doesn't depend on n .

Recall

The method converges on $[t_0, T]$ if

$$\lim_{N \rightarrow \infty} \left(\max \{ \|e_n\| : n = 1, \dots, N \} \right) = 0$$

and we are trying to show Euler's method converges...

$$\lim_{N \rightarrow \infty} \max \left\{ \begin{array}{l} \text{propagated error} \\ e^{(T - t_0)\lambda} \|e_0\| + \frac{ch}{\lambda} [e^{(T - t_0)\lambda} - 1] \end{array} \right.$$

this term goes to zero in the limit.
 generated error...
 generated by the truncation in each step of the integral on the right

$$h = \frac{T - t_0}{N} \rightarrow 0 \quad \text{as } N \rightarrow \infty$$

$$= e^{(T - t_0)\lambda} \|e_0\| = 0$$

up to the rounding error in the initial condition...

This also tells us the rates of convergence when there is no rounding error...

$$\text{total error} \leq \frac{Ch}{\lambda} \left[e^{(T-t_0)\lambda} - 1 \right] = O(h)$$

This is called a first order method...

Generalize the argument...

$$y' = f(t, y)$$

$$y(t_0) = y_0$$

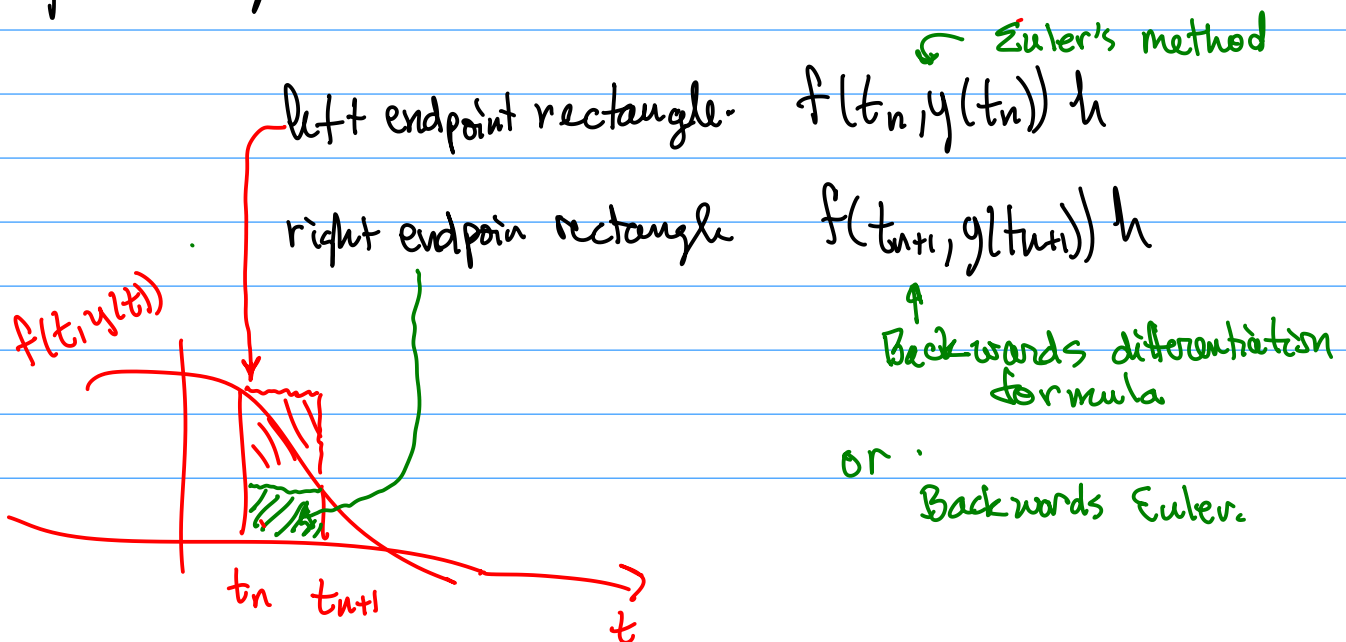
$$t_n = t_0 + hn$$

$$h = \frac{T - t_0}{N}$$

Integrate the ODE on both sides from t_n to t_{n+1}

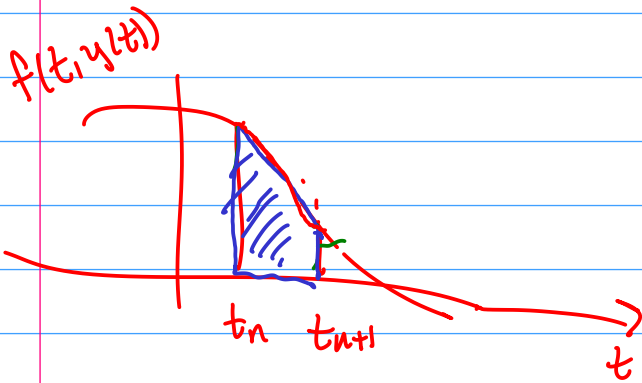
$$\int_{t_n}^{t_{n+1}} y'(t) dt = \int_{t_n}^{t_{n+1}} f(t, y(t)) dt$$

$$y(t_{n+1}) - y(t_n) = \text{Some approximation...}$$



Trapezoid rule

$$\frac{f(t_n, y(t_n)) + f(t_{n+1}, y(t_{n+1}))}{2} h$$



Introduce a parameter θ and approximate the integral as a weighted average of the left and right endpoints

Theta Method

$$\left[\theta f(t_n, y(t_n)) + (1-\theta) f(t_{n+1}, y(t_{n+1})) \right] h$$

Two questions:

- ① Does it converges as $N \rightarrow \infty$?
- ② How fast?

Theoretical assumption of no rounding error.

Trapezoid rule:

$$y(t_{n+1}) - y(t_n) \approx \frac{f(t_n, y(t_n)) + f(t_{n+1}, y(t_{n+1}))}{2} h$$

Numerical Scheme

$$y_{n+1} = y_n + \frac{f(t_n, y_n) + f(t_{n+1}, y_{n+1})}{2} h$$

use this equation to solve for y_{n+1} given y_n .

Since y_{n+1} appears on the right side in the function f , this is called an implicit method.

Solving for y_{n+1} using Newton's method is not so bad because you can use y_n as the initial guess so it only takes 1 or 2 iterations to find y_{n+1} .

computer lab on practical implicit methods.

Analysis of this:

$$y_{n+1} = y_n + \frac{f(t_n, y_n) + f(t_{n+1}, y_{n+1})}{2} h$$

Truncation error τ_n . plug the exact solution into the numerical scheme and compute the residual error.

$$y(t_{n+1}) - y(t_n) = \frac{f(t_n, y(t_n)) + f(t_{n+1}, y(t_{n+1}))}{2} h + \tau_n$$

Estimate τ_n :

$$\tau_n = y(t_{n+1}) - y(t_n) - \frac{f(t_n, y(t_n)) + f(t_{n+1}, y(t_{n+1}))}{2} h$$

mult by h here

Since $y' = f(t, y)$ then $f(t_n, y(t_n)) = y'(t_n)$

and $f(t_{n+1}, y(t_{n+1})) = y'(t_{n+1})$

Taylor series expand everything

Set $t = t_n$ for convenience

$$y(t_{n+1}) = y(t+h) = y(t) + h y'(t) + \frac{h^2}{2} y''(t) + O(h^3)$$

$$y'(t_{n+1}) = y'(t+h) = y'(t) + h y''(t) + O(h^2)$$

In the end everything cancels, check this on Thursday...

$$\tau_n = O(h^3).$$